## Hilbert Sort

Sorting numerical data not only makes it easy to search for a particular item, but also makes better use of a CPU's cache: any segment of data that's contiguous in memory will describe a set of items that are similar in some sense. Things get more complicated if our data represents points on a 2D grid. If points $(\boldsymbol{x}, \boldsymbol{y})$ are sorted by $\boldsymbol{x}$, breaking ties by $\boldsymbol{y}$, then adjacent points will have similar $\boldsymbol{x}$ coordinates but not necessarily similar $\boldsymbol{y}$, potentially making them far apart. To better preserve distances, we can sort the data along a space-filling curve.

The Hilbert curve starts at the origin ( $\mathbf{0}, \mathbf{0}$ ), finishes at $(\mathbf{s}, \mathbf{0})$, in the process traversing every point in axis-aligned square with corners at ( $\mathbf{0}, \mathbf{0}$ ) and ( $\mathbf{s}, \mathbf{s}$ ). It has the following recursive construction: split the square into four quadrants meeting at ( $s / \mathbf{2}, \boldsymbol{s} / \mathbf{2}$ ). Number them 1 to 4 , starting at the lower left and moving clockwise. Recursively fill each of them with a suitably rotated and scaled copy of the full Hilbert curve.

Start with a single point at $(\boldsymbol{s} / \mathbf{2}, \mathbf{s} / \mathbf{2})$. Then, repeat these steps:

- Scale and copy the current construction into each of the 4 quadrants.
- Rotate quadrant 1 by -90 degrees and flip it vertically, so that the start of the curve is closest to the lower left corner ( $\mathbf{0 , 0} \mathbf{0}$ ).
- Rotate quadrant 4 by 90 degrees and flip it vertically, so that the end of the curve is closest to the lower right corner ( $\mathbf{s}, \mathbf{0}$ ).
- Now, connect the end of the curve in quadrant 1 to the start of the curve in quadrant 2 , connect the end of quadrant 2 to the start of quadrant 3 , and the end of quadrant 3 to the start of quadrant 4.

Here are the first two iterations:



The Hilbert Curve is built by repeating this construction infinitely many times. The following diagram shows the first six steps of building the Hilbert Curve:


Given some places of interest inside of a square region, sort them according to when the Hilbert curve visits them, starting from ( 0,0 ). Without going into gory detail about Fractal theory, note that making $\boldsymbol{s}$ odd guarantees that all integer points are visited just once, so their visitation order in relation to each other is unambiguous.

## Input

Each input will consist of a single test case. Note that your program may be run multiple times on different inputs. The first line of input contains two space-separated integers $n$ and $\boldsymbol{s}\left(1 \leq \boldsymbol{n} \leq 100,000,1 \leq \boldsymbol{s}<10^{9}, \boldsymbol{s}\right.$ is odd). The next $\boldsymbol{n}$ lines describe locations of interest by space-separated integers $\boldsymbol{x}$ and $\boldsymbol{y}(0 \leq \boldsymbol{x}, \boldsymbol{y} \leq s)$. No two locations will share the same position.

## Output

Output the $\boldsymbol{n}$ ordered pairs, one per line, with $\boldsymbol{x}$ and $\boldsymbol{y}$ separated by a space, Hilbertsorted according to their positions.

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| Sample Input | Sample Output |
| :---: | :---: |
| 1425 | 55 |
| 55 | 105 |
| 510 | 1010 |
| 520 | 510 |
| 105 | 520 |
| 1010 | 1020 |
| 1015 | 1015 |
| 1020 | 1515 |
| 155 | 1520 |
| 1510 | 2020 |
| 1515 | 2010 |
| 1520 | 1510 |
| 205 | 155 |
| 2010 | 205 |
| 2020 |  |

